Goal Decomposition Tree: an agent model to generate a validated agent behaviour

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Abstract. This paper deals with a goal-oriented agent model called Goal Decomposition Tree (GDT) allowing both to specify and validate the behaviour of an agent. This work takes place in a global framework whose goal is to define a process allowing to start from a problem specification to obtain a validated implementation of a corresponding MAS. The GDT model has been used to specify a prey-predator system which has been verified this way.

1 Introduction

Our research deals with methods and models, in order to help multiagent systems designers to manage the complexity of MAS. We also aim at helping to develop multiagent systems whose behaviour can be verified. Our proposal takes place in "formal transformation systems" context as defined in [3]: "Formal transformation systems provide automated support to system development, giving the designer much more confidence that the resulting system will operate correctly, despite its complexity". The main principle of a formal transformation system is that each transformation step must preserve correctness from one model to the next. This constraint has been used as a guide to design our approach which consists of four steps:

- 1. **an agentification method** that helps to determine the set of agents which must be used to implement a given system;
- 2. an agent design model to help to design an agent behaviour that can be verified:
- 3. a **proof system** to prove that the agent model satisfies the main goal of the agent:
- 4. **an implementation model** that can be automatically generated from the agent design model.

Our aim is to provide a complete MAS design process starting from the problem specification and ending by a MAS implementation. Moreover, as it is said in [3] about formal transformation systems, "since each transformation preserves correctness from one model to the next, the developer has much more confidence that no inconsistencies or errors occurred during the design process". Our agentification method and our implementation model have already been presented in other articles [15,11]. This paper is focused on the agent model called GDT.

In order to be able to implement and validate an agent, its behaviour must be clearly and formally specified. In our proposal, this specification is based on a Goal Decomposition Tree (GDT) which helps to describe how an agent can manage its goals. It is used

- as a support for agent behaviour validation and proof;
- to guide the implementation of the agent behaviour induced by the tree using an automaton as in our implementation model called SPACE [11].

The main contribution of this model is that it can be proven that the specified agent behaviour is correct according to the main goal of this agent. The proof mechanisms are detailed in [4]. That's why the aim of this model is to provide a declarative description of goals. Several works have already pointed out the advantage to have a declarative description of goals [19], [6], [17], [12], [5]. Many multiagent models or systems are essentially focused on procedural aspects of goals which is important in order to obtain an executable agent. But the declarative aspect of goals is also very important. Indeed, as it is said in [19], "by omitting the declarative aspect of goals, the ability to reason about goals is lost. Without knowing what a goal is trying to achieve, one can not check whether the goal has been achieved, check whether the goal is impossible". In [17], the authors say that declarative goals "provide for the possibility to decouple plan execution and goal achievement". A GDT is a partial answer to these requirements: as it will be shown in next sections, both procedural and declarative aspects of goals management can be described by a GDT.

Another important aspect of a GDT is that it is intended to be used to directly generate the behaviour of the agent to which it is associated. Indeed, as explained in [17], "the biggest role of goals in agents is thus not the ability to reason about them but their motivational power of generating behaviour". Moreover, "A certain plan or behaviour is generated because of a goal". It is exactly what a GDT allows to express. Nodes of a GDT are goals the agent has to achieve. As in [19], goals are considered as states of the world the agent has to reach.

Inside the tree, a goal is decomposed into subgoals using decomposition operators. The notion of subgoal used here is similar to the one described in [17]: "a goal can be viewed as a subgoal if its achievement brings the agent closer to its top goal". The notion of "closeness" used by van Riemsdjik et al. is specified differently by each decomposition operator. In the same paper, authors distinguish "subgoals as being the parts of which a top goal is composed" and "subgoals as landmarks or states that should be achieved on the road to achieving a top goal". In a GDT, these two kinds of subgoals exist. The fact that a subgoal is of a particular kind is a direct consequence of the decomposition operator which has been used to introduce the subgoal. All works on agent goals management do not use the same notion of subgoal. In [17] or [6], a subgoal is considered as a logical consequence of a goal. So, in these works, subgoals can be considered as necessary conditions for the satisfaction of the parent goal. In our vision, subgoals are on the contrary sufficient conditions for the satisfaction of the parent goal. The TAEMS model [18] does not use goals buts tasks which, however, can be compared to goals used in our work.

A decomposition operator encapsulates a set of mechanisms corresponding to a typical goals management behaviour ([6], [17], [19]). Each operator is specified by different kinds of semantics:

- a goal decomposition semantics describing how a goal can be decomposed into subgoals with this operator;
- a semantics describing how to deduce the "type" of the parent goal knowing the types of its subgoals;
- a semantics associating an automata composition pattern to each operator.
 These patterns are used incrementally to build the complete automaton describing the agent behaviour;
- a semantics associating a local proof schema. This schema is used to verify the agent behaviour (i.e. to prove that its goals management behaviour satisfies its main goal). This semantics is described in [4].

Section 2 defines the notion of goal as it is used in this work and describes the typology of goals which has been defined. Section 3 describes the set of operators which can be used to decompose a goal into subgoals inside a GDT. For each operator, the two first semantics described before are given. Section 4 defines more precisely a GDT and shows how it can be built using the tools described in the two previous sections. Last but not least, section 5 presents a comparison of the proposed model with other works.

2 Goals and typology of goals

In the context of a GDT, a name and a satisfaction condition are associated to each goal. According to the type of each goal, additional information is also used to completely specify the goal. Satisfaction conditions are used to specify goals formally with respect to the declarative requirement for goals described in the previous section. A goal is considered to be achieved if its satisfaction condition is logically true. Satisfaction conditions are expressed using a temporal logic formalism which is a subset of TLA [10]. More precisely, primed variables have been used. For example, if x is a variable, the notation x in a satisfaction condition corresponds to the value of x before the execution of the goal to which the condition is associated. On the contrary, the notation x' corresponds to the value of x after the goal execution. Note that more complex temporal logic formulae are also used to specify the semantics of decomposition operators described in section 3. These formulae, which are not described in the paper because of a lack of space, are then used during the proof process. However, an example of such a formula is given in section 3.1 for the sequand operator.

Variables are supposed to be attributes managed by the agent. Thus, specifying goals of an agent helps also to define the set of variables defining the view of the agent on its environment. For example, in [1], the authors describe a case study where two robots are collecting garbage on Mars. One of the two robots, named R1, must move in order to discover pieces of garbage which must be picked up. As a consequence, for R1, to be in a location where there is a piece of garbage corresponds to a goal. The satisfaction condition of this goal can be defined by: garbage = true where garbage is supposed to be an internal

variable of the agent which describes its perception of its current location. This variable is supposed to be updated each time the agent moves i.e. each time its coordinates (x,y) (which are also internal variables of the agent) are modified.

A typology of goals has been defined in order to distinguish more precisely different ways to manage goals decomposition. The type of a goal has consequences on the execution process of this goal, on the design of its corresponding behaviour automaton and on the proof process of the behaviour implied by its execution by the agent.

The first criterion to distinguish goals corresponds to the situation of the goal in the tree. This leads naturally to distinguish two first kinds of goals: elementary and intermediate goals.

Elementary goals: they correspond to the leaves of the tree that's why they are not decomposed into subgoals. Furthermore, they are not only defined by a name and a satisfaction condition but also by a set of actions. The execution of these actions are supposed to achieve the goal i.e. to make its satisfaction condition true. Notice that these actions are related to the capabilities of the agent as described in [6]. They correspond to the procedural aspect of goals described in the previous section. As satisfaction conditions, they are based on variables of the agents. The aim of an action is to modify these variables. For example, for the robot R1 described in section 2, moving one step right can be an elementary goal. Its satisfaction condition is: $x' = x + 1 \land y' = y$. The corresponding actions are: x := x + 1; yarbage := containsGarbage(x, y). It is supposed that yarbage := containsGarbage(x, y). It is supposed that yarbage := containsGarbage(x, y).

Intermediate goals: They correspond to the nodes of the tree which are not leaf ones. They are specified by a name, a satisfaction condition and also a Local Decomposition Tree (LDT). A LDT contains a root node corresponding to the intermediate goal and a decomposition operator which creates as many branches (and subgoals) as needed by the operator. It describes how the intermediate goal can be decomposed into subgoals, and sometimes in which context, in order to be achieved. The number and the order of subgoals to be achieved in order to achieve the parent goal depends on the chosen operator (see next section for more details).

The second criterion used to define goals is related to the goals satisfiability. Using this criterion, two kinds of goals are again distinguished: Necessarily Satisfiable goals (NS) and Not Necessarily Satisfiable goals (NNS).

Necessarily Satisfiable goals (NS): This kind of goals ensures that, once the goal has been executed, the satisfaction condition of the goal is necessarily true (the goal is achieved).

Not Necessarily Satisfiable goals (NNS): this set of goals is complementary to the previous one. It is the more prevalent case. This kind of goals is not necessarily achieved after its actions or its decomposition (and the subgoals associated to this decomposition) have been executed. This kind takes into account that some actions or some decompositions can only be used in certain execution contexts.

If these contexts are not set when the goal is to be executed, these actions or decompositions become inoperative.

This criterion is used for the automaton generation step and the proof process. It is not involved in the execution process of a goal. Moreover, during the GDT design step, using this criterion may help to verify that the tree is syntactically correct. Indeed, as it will be shown in the next section, decomposition operators do not accept all kinds of goals from this criterion point of view.

The third and last criterion used to distinguish goals is related to the evaluation of the satisfaction condition. Using this criterion, two other kinds of goals can be defined: Lazy goals (L) and Not Lazy goals (NL).

Lazy goals (L): when a goal is considered to be a lazy one, its satisfaction condition is evaluated before considering its actions (for an elementary goal) or its decomposition (for an intermediate one). If the satisfaction condition is true, the goal is considered to be achieved which implies that the set of actions or the decomposition associated to the goal are not executed. If the satisfaction condition is false, the set of actions or the decomposition are executed. As a consequence, primed variables can not be used in the satisfaction condition associated to a lazy goal. Indeed, in this case, implicit stuttering of variables is supposed. It means that if a variable is not used in the satisfaction condition, its value is considered as unchanged.

Not Lazy goals (NL): this set of goals is complementary to the previous one. For a not lazy goal, the associated set of actions or decomposition tree is always executed even if the satisfaction condition is already true before the execution. This property is usefull for goals which must absolutely always be executed, whatever is the system execution context.

This criterion can be compared to the requirement for goals management described in [19]: "...given a goal to achieve condition s using a set of procedures (or recipes or plans) P, if s holds, then P should not be executed". This requirement corresponds to the solving process of a lazy goal whose satisfaction condition is already satisfied at the beginning of the process.

As a conclusion, in the context of a GDT, each goal can be characterised by three criteria which can be combined independently. Figure 1 summarizes the graphical notations introduced for the two last criteria. Each criterion has two possible values which implies that eight effective kinds of goals can be used in the tree.



Fig. 1. NS, NNS, L and NL goals

Formally, a goal is described by a 6-tuple < name, sc, el, ns, lazy, LDT or actions > with: name a string, sc (satisfaction condition) a temporal logic formula, el (elementary) a boolean, ns (necessarily satisfiable) a boolean, lazy a boolean and a Local Decomposition Tree (LDT) or a set of actions (actions).

From this definition, the process of goal execution in the GDT can be described by the algorithm given in figure 2. The *solve* function used in this algorithm describes how the tree must be walked during the execution process.

```
boolean solve(G) :
   if (G.lazy) then
      if (G.sc) then return(true);
      endif
   endif
   if (G.el) then execute(G.actions);
   else satisfy(G.LDT);
   endif
   return(G.sc);
```

Fig. 2. Algorithm for executing a goal in a GDT

The *execute* function consists in executing sequentially the set of actions associated to the goal. The *satisfy* function is a recursive one and uses itself the *solve* function which just has been defined. The *satisfy* function is detailed in the next section which also describes available decomposition operators.

3 Decomposition operators

In this section, available decomposition operators are described. For each operator, the two first semantics are given that is to say the decomposition semantics and the goals types composition semantics. The goals types composition semantics is based only on the satisfiability mode of goals (NS or NNS). Indeed, the other criteria have not a direct influence on decomposition operators. AND and OR operators are available but are not described because they correspond to the standard logical operators. In a tree, they are managed as indeterminist operators.

Before describing each operator, let us precisely define goal decomposition. Let A be the goal to be decomposed and Op the decomposition operator to be used. As almost all available operators are binary, Op is also supposed to be binary (but it does not modify the semantics of the decomposition). Let B and C be the subgoals introduced by the use of Op to decompose A. As described in the previous section, Op(B,C) corresponds to the Local Decomposition Tree associated to A. The semantics of this decomposition is that the satisfaction of Op(B,C) (i.e. the LDT) implies the satisfaction of A. But what does mean the satisfaction of Op(B,C)? It corresponds exactly to the decomposition semantics of each operator. Indeed this semantics describes how many subgoals must be achieved and in which order to be able to satisfy the parent goal. In other words, the satisfy function used inside the solve function given previously is different

for each operator. So in the sequel, for each operator, the *satisfy* function is instantiated. Let us notice that, for all operators, this function uses the *solve* function in order to evaluate the satisfaction of subgoals.

3.1 SeqAnd operator

This operator corresponds to a "sequential And" operator. Indeed, the main difference with the standard logical And operator is that the two subgoals must be executed in the order specified by the operator. The figure 3(a) gives the satisfy function corresponding to this behaviour. The composition semantics is the same as the And operator's one i.e. the parent goal is NNS unless the two subgoals are NS.

```
boolean satisfy(SeqAnd(B,C)):
if (solve(B))
then return (solve(C));
else return(false);
endif

(a) SeqAnd

(b) SeqOr
```

Fig. 3. SeqAnd and SeqOr satisfaction algorithms

Last but not least, here is the temporal logic formula associated to this operator. This formula is supposed to be true when decomposing a lazy parent goal A in two subgoals B and C using seqand operator. It is a logical expression of the "solve" function with G being A, G.ldt being (B seqand C) and G being a lazy goal. There exists a second formula to be used when the parent goal is a not lazy one. This formula is used during the proof process:

$$\Box(\Diamond \neg A \land (\neg A \Rightarrow \Diamond B) \land ((\neg A \Rightarrow \Diamond B) \Rightarrow \Diamond C) \Rightarrow \Diamond A)$$

 \Box is a temporal logical operator meaning "always". \Diamond is an other temporal logical operator meaning "eventually". Informally, this formula means that if A is not satisfied (i.e. its satisfaction condition is false; this condition must be verified because A is a lazy goal), and if B can be satisfied when A is not satisfied and if C can be satisfied in the previous context (i.e. B can be satisfied when A is not satisfied) then A will eventually be satisfied (i.e. its satisfaction condition will eventually be true).

3.2 SegOr operator

The difference between SeqOr and the standard logical Or operator is the same as the one between SeqAnd and the standard logical And operator described in the previous section. The figure 3(b) gives the satisfy function associated to SeqOr. Its composition semantics is the same as the Or operator's one i.e. the parent goal is NNS unless the two subgoals are NS.

3.3 SyncSeqAnd operator

This operator is a synchronized version of the SeqAnd operator. Unlike SeqAnd, this operator ensures that the two subgoals (if they are both executed) are

solved without any interruption by another agent. This operator must not be used too much. By reducing the shared vairables, the agentification method we have proposed [7] is designed to limit cases where this kind of operators must be used. Its decomposition semantics, goals types composition semantics and the corresponding satisfaction algorithm are the same as the SeqAnd operator's ones.

3.4 SyncSeqOr operator

The difference between SyncSeqOr and SeqOr is the same as the one between SyncSeqAnd and SeqAnd described in the previous section. This operator is a synchronized version of the SeqOr operator.

3.5 Case operator

This operator decomposes a goal into subgoals according to conditions defined by logical formulae. These logical formulae use the same variables as satisfaction conditions. The decomposition semantics of this operator states that the disjunction of conditions must always be true when the parent goal is decomposed. The principle is that if a condition is true, the corresponding subgoal must be executed. The satisfaction of the parent goal depends on the satisfaction of the chosen subgoal. This semantics is summarised by the associated satisfy function given in figure 4.

```
boolean satisfy(Case(B1,B2,cond_B1,cond_B2)):
if (cond_B1)
then return(solve(B1))
else if (cond_B2) then return(solve(B2))
```

Fig. 4. Case satisfaction algorithm

As far as the composition semantics of the operator is concerned, there are four possible trees as shown in figure 5. If subgoals are both necessarily satisfiable, the parent goal is necessarily satisfiable. If at least one of the subgoals is not necessarily satisfiable, the parent goal is not necessarily satisfiable. It is very important to notice that the property of being "necessarily satisfiable" is a little bit different in the context of the case operator. Indeed, here, a subgoal is necessarily satisfiable only if its associated condition is true. For the other operators, when a goal is declared to be necessarily satisfiable, it is true in any context. This characteristic is particularly useful for the proof process.

3.6 Iter operator

This operator is an unary one. The main feature of this operator is that its behaviour depends on the satisfaction condition of the parent goal. The decomposition semantics of this operator states that the parent goal will be achieved if the subgoal is achieved several times. In other words, the satisfaction condition of the subgoal must be true several times in order the satisfaction condition of the parent goal to become true. Nevertheless, it is possible that, for example,

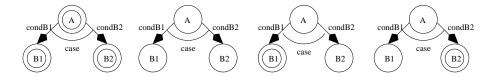


Fig. 5. Case operator composition semantics

modifications in the environment of the agent imply that the satisfaction condition of the parent goal becomes true even if the subgoal has not really been satisfied. In that case, the parent goal is considered to be satisfied and the solving process of the subgoal must stop.

This operator is very important because it takes into account a progress notion inside a goal execution process. For example, let suppose that the satisfaction condition of the parent goal A is "to be in (x,y) location". Let suppose that the agent can only move one step at a time. As a consequence, the execution of A must be decomposed into n execution of the subgoal "move one step", n being the number of steps between the current location of the agent and the desired final location.

This operator can only be used when the satisfaction of the subgoal implies a progress in the satisfaction of the parent goal. In other words, each time the subgoal is satisfied, the satisfaction of the parent must be closer. However, sometimes it is possible that the subgoal can not be satisfied (because the context of the agent has changed for example). In this case, the satisfaction degree of the parent goal stays at the same level and the subgoal must be executed again. The important characteristic of this operator is that the satisfaction level of the parent goal can not regress after a satisfaction step of the subgoal, even if this step has failed. If it is the case, it means that the *Iter* operator should not have been used. The proof schema associated to the *Iter* operator helps to verify this property.

The overall behaviour of the operator described in the previous paragraph is summarised by the associated satisfy function given in figure 6(a).

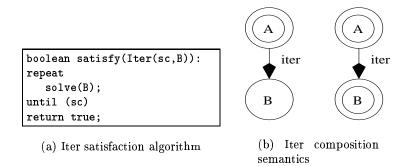


Fig. 6. Iter Operator

The goals types composition semantics of this operator is summarised in figure 6(b). It shows that the subgoal can be either necessarily satisfiable or not. However, the parent goal is always necessarily satisfiable. Indeed, the behaviour of the operator implies that the solving process of the subgoal stops when the satisfaction condition of the parent goal is true which implies that this one is necessarily satisfiable.

3.7 Comparison with other works on agent goals decomposition

Other works on goals management of an agent propose mechanisms to express relations between goals or subgoals. In this paragraph, two of them are detailed in order to be precisely compared with GDT. In GOAL [6], the authors propose a global logical framework in order to formalise the goals management behaviour of an agent. In this framework, the state of an agent is defined by a mental state $\langle B,G \rangle$ which consists of the beliefs and goals of the agent. Beliefs and goals are modelled by logical formulae: B(F) is a belief and G(F) is a goal, F being a logical formula.

In this framework, a goal must not be deductible from the set of beliefs. When performing an action, if a goal becomes a consequence of the modifed set of beliefs, it is removed from the goals base. The goal is considered to be achieved.

The behaviour of the agent is specified by a set of *conditional actions*. A conditional action is a pair $\phi \to do(a)$ where ϕ is a condition (a logical formula) and a is an *action*. There are three kinds of actions:

- beliefs management actions: theses actions allow to manage the set of beliefs:
 - $ins(\phi)$ adds $B(\phi)$ to the set of beliefs,
 - $del(\phi)$ deletes $B(\phi)$ from the set of beliefs;
- goal management actions: these actions allow to explicitly manage goals:
 - $adopt(\phi)$ adds $G(\phi)$ to the set of goals,
 - $drop(\phi)$ deletes $G(\phi)$ from the set of goals;
- basic actions: these actions are described by a semantic function \mathcal{T} , a partial function that modifies the set of beliefs of the agent.

For instance, here is how our SeqAnd operator could be translated in Goal. Let suppose that A is a goal that is decomposed in SeqAnd(X,Y), that is to say $SeqAnd(X,Y) \Rightarrow A$.

The behaviour of the agent corresponding to the execution of the goal A can then be described by the following conditional actions:

- $-G(A) \wedge \neg B(X) \rightarrow do(adopt(X))$ (if A must be satisfied and X is not yet believed, then X becomes a goal of the agent);
- $G(A) \wedge B(X) \wedge \neg B(Y) \rightarrow do(adopt(Y))$ (if A must be satisfied and X has already been achieved (and is thus a belief), then Y becomes a goal of the agent);
- $-G(A) \wedge B(X) \wedge B(Y) \rightarrow do(ins(A))$ (if A must be satisfied and X and Y have been achieved, then A is achieved and can be added to the set of beliefs of the agent. It will also be removed from the set of goals of the agent because Goal agents implement the blind commitment strategy).

It is assumed that there are also rules to satisfy goals X and Y which are supposed to add B(X) and B(Y) to the set of beliefs. However, with our model, X and Y are removed from the set of goals remaining to solve by the agent after the resolution of A. This can not be expressed in GOAL because conditional actions can not be combined, for instance to be sequentialised. More generally, the hierarchical structure of our model allows a progressive specification of the agent behaviour which is more difficult with Goal. Moreover, more elements can be proven with our model than with GOAL. For example, relations between goals like "ITER" can not be proven with GOAL. Last but not least, our model allows to perform proofs using first order logic which is not the case with GOAL.

Our decomposition operators can also be compared to the Quality Accumulation Functions (QAF) proposed in TAEMS [18]. TAEMS is a modelling language allowing to describe activities of agents operating in environments where responses by specific deadlines may be required. That's why TAEMS represents agent activities in terms of a task structure i.e. a graph where tasks can be decomposed into subtasks using QAFs. A QAF specifies how the quality of a task can be computed using qualities of its substasks. The quality of a task evaluates the level of completion of the task. QAF can be seen as a quantitative version of our logical decomposition operators.

4 The GDT design process

A Goal Decomposition Tree (GDT) specifies how each goal can be achieved by an agent. More precisely, the root node of the tree is associated to the main goal of the agent, i.e. the one which is assigned to the agent by the used agentification method ([15] [20]). If this goal is achieved by the agent, the agent is considered to be satisfied from the multiagent system point of view. The tree describes how this goal can be decomposed in order to be achieved using a solution which must be the most adapted to the agent context as possible. Notice that the overall tree can be seen as a collection of local plans allowing to achieve each goal. A local plan corresponds to a Local Decomposition Tree associated to a subgoal. The main difference with plans used in [1] is that, in a GDT, they are organised hierarchically. A GDT is very close to the tasks graph used in TAEMS [18]. This graph describes relationships between tasks agents may have to achieve. A graph, instead of a tree, is needed because relationships between goals ("disables", "facilitates"...), different from decomposition ones, can be expressed. These relationships often involve tasks which are executed by different agents. That is why they can not be expressed in a GDT but are taken into account at the system level (which is not presented here).

The building process of a GDT consists of four steps. In a first step, a tree must be built by the designer, starting from the main goal of the agent using a top-down process. This first step allows to introduce subgoals with their satisfaction condition, elementary goals with their associated actions and also decomposition operators. The designer must also decide for each goal if it is lazy or not. During this step, the designer must also define invariants associated to

the tree. A GDT must be designed so as to ensure that the execution of goals is always consistent with all the invariants. This characteristic is verified during the proof process. Invariants specify properties of the system inside which agents are to be executed. For example, for a prey/predator system, invariants must specify that only one agent can be on a grid cell at a time.

In order to make the building process of the tree easier, we are currently defining what can be seen as design patterns i.e. rules which can be used to choose the right operator in particular contexts. For example a rule is focused on the problem of interdependency between goals. When this property exists between two goals A and B, it means that the satisfaction of A has an influence on the satisfaction of B. When detected, this property can help to guide the choice of the decomposition operator. For example, let suppose that a goal G can be satisfied if two subgoals B and C are satisfied. The And operator may be used to model this hypothesis. But if another hypothesis indicates that the satisfaction of B can prevent from the satisfaction of C, the And operator can not be used anymore, but must be replaced by the SeqAnd operator.

In the second step of the GDT design process, the designer must decide for each elementary goal if it is necessarily satisfiable or not. In a third step, the type of each intermediate goal, as far as satisfiability is concerned, is computed using the goals types composition semantics of each used decomposition operator. Unlike the first step, this step is a down-top process. During this process, inconsistencies or potential simplifications can be detected. In that case, the first step must be executed again in order to modify the tree. Figure 7 shows such a simplification for the SeqOr operator which can be detected during this step. The first tree can be replaced by the second tree because if the first subgoal of a SeqOr is a necessarily satisfiable one, the second subgoal will never be executed.

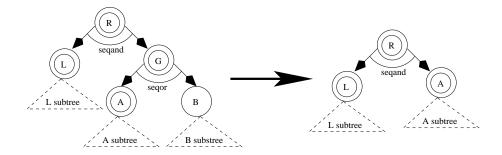


Fig. 7. Tree simplification with SeqOr

Once the three first steps have been achieved, a proof of the tree can be built in a fourth step. The process used to achieve this proof is described in [4]. Again, this step can lead to detect inconsistencies in the tree based on proof failures. In a last step, the validated tree is used to build the behaviour automaton of the agent which can then be implemented. This process is also described in [16].

As explained before, the building of the tree leads also to the definition of variables and actions of the agent which are essential parts of an agent model. As a consequence, the GDT and the associated design process can be seen as a design tool for a validated agent model in the context of a MAS design.

5 Comparison with other models and methods

5.1 Comparison with other goal-oriented agent models

Model	Goal	Goal	Action kinds	Plan
	Expression	management		language
		hypotheses		(operators)
Winikoff et	satisfaction	DS, DI, CG,	GA, BI, BD, BAN	sequencing, paral-
al.	failure	PG, KG		lelism, conditional
				selection
AgentSpeak	none	none	GA, BI, BD, BAN	and, conditional
				selection (context)
Goal	satisfaction	DS	GA, GD, BI, BD, BAS	only atomic condi-
				tional actions
GDT	satisfaction	DS	GA, GD, BI, BD de-	many
			rived from the GDT,	
			BAS	

Table 1. Comparison with goal-oriented agent models

Table 1 compares our agent model with a few other ones with a goal oriented point of view: Winikoff et al's model [19], AgentSpeak [14] and GOAL [6]. In the Goal expression column, it is specified whether a formal satisfaction condition and a formal failure condition is expressed for each goal in the model. For the models having only a procedural point of view, like AgentSpeak, there is no formal expression of goals. Only the Winikoff et al's model explicitly gives a formal failure condition, making a distinction between a plan failure and a goal failure.

Among the Goal management hypotheses, we distinguish the five characteristics described in [19, 13]. The Drop Successful attitude (DS) consists in dropping a goal that is satisfied. The Drop Impossible attitude (DI) consists in dropping a goal that becomes impossible to satisfy. A goal is persistent (PG) if it is dropped only if it has been achieved or if it becomes impossible to achieve. The other characteristics correspond to constraints on the goal management process. The Consistent Goals property (CG) is satisfied if the set of goals the agent has to achieve must be consistent (if a is a goal, not(a) cannot be a goal). Finally, the Known Goals (KG) property specifies that the agent must know the set of all its goals. DS property implicitly holds in our model: it is a direct consequence of the execution schema of a GDT specified by the solve and satisfy functions described previously. DI does not hold because, as there is not any failure condition associated to goals, there is no way to decide whether a goal become impossible

to achieve. PG does not hold because DI does not hold. Last but not least, CG and KG properties which are constraints on the model usage, do not hold in our model because it does not need these two properties to be verified in order to be used.

In the *Action kinds* column, kinds of actions the language provides are precised. These actions can be divided into 3 types: actions concerning goals management (goal dropping GD, goal adoption GA), actions concerning beliefs management (belief insertion BI, belief deletion BD) and all other actions that we call Basic Actions. These actions may be *Specified* in the language (BAS) or only Named (BAN). BAS are essential to allow a proof process.

Finally, in the last column, we tried to enumerate the *Goal decomposition operators* provided by the language. For the model described in this paper, see section 3. The plan language of Goal is rather a rule language. But for each rule, only one action may be executed: there is, for instance, no sequence operator. In a GDT, plans rely on goal decompositions, and as a consequence, the expressivity of our plan language is also the expressivity of our goal decomposition language.

5.2 Comparison with goal-oriented MAS development methods

Already existing MAS development methodologies like Tropos [2], Prometheus [12], Kaos [5] or MaSE [3] include goals management. As these methodologies, our approach is intended to be a start to end goal-oriented support for MAS design with a specific property: to be able to prove that the generated behaviour is correct with respect to system requirements. As previously explained, our approach is very close to MaSE from this point of view: they are both designed as "Formal Transformation Systems".

All these methods lead to define system goals which are then used as a base for further design steps. But, as explained in [9], "beyond the initial phases of requirements, the goal-oriented focus of these methodologies dissipates and does not progress into the detailed design phase where agent internals are fleshed out". The work presented in this paper tries exactly to answer to this problem because it is focused on agents internal from a goal oriented point of view. However Tropos and a new version of Prometheus presented in [9] try to keep goals during all the design process.

In the following, previous cited methodologies are compared from a goals management point of view (for the Prometheus methodology, the comparison is based on [9]). Used criteria are the following:

- Goals level: goals can be defined at the system level(SL), at agents level(AL)
 or at agents'internal level (AIL). It is supposed that AIL implies AL and that
 AL implies SL.
- Goals description : goals can be specified formally (SF) or are just named (N).
- Goals role: goals are directly exploited to produce models during the design phase (PM) or are used to ensure that systems requirements are taken into account (SR).

- Typology use: when a typology exists, it can be used to help to understand goals (UG) or to help to determine goals (DG) or to guide strongly further design steps (GD).
- Goals decomposition : does the method include means to express relations between goals and subgoals ?

	\mathbf{GDT}	Prometheus	MaSE	KAOS	Tropos
Goals level	AIL	AIL	AL	SL	AIL
Goals description	SF	N	N	SF	N
Goals role	PM	SR	SR	SR	PM
Typology use	GD		DG	UG	UG
Goals Decomposition	operators	$\mathrm{and/or}$	typology	and/or	${ m and/or/contribution}$

Table 2. Comparison with Goal-Oriented MAS development methods

As far as goals decomposition is concerned, the term "typology" used for Mase means that available decompositions are directly defined by the different types of goals.

6 Conclusion

In this article, we presented a goal-oriented behaviour model of an agent relying on a Goal Decomposition Tree. The *qoal* notion is central in the development of an agent. This appears for instance in the desires concept of the BDI model, or is the basis of other methods such as Goal or Moise ([6], [8]). Using a decomposition tree, the user can progressively specify the behaviour of an agent. Thus, goals and plans are closely linked: the decomposition tree of a goal is the plan associated to this goal. A part of goal decomposition operators involves nondeterminism, which is necessary for autonomous agents. Of course, using our model, an agent designer must specify each goal by a satisfaction condition. This may seem difficult, but the experience shows that rapidly, inexperienced designers can write the right satisfaction condition. Moreover, this model can be verified using our proof method. The model can then be automatically translated into a behaviour automaton which is, as a consequence, also validated. This automaton can be implemented inside agents which can be developed using any MAS development platform in order to manage their life cycle, their communications, etc. However, the design and the proof process are strongly disconnected. So, the designer can develop the GDT without taking care of the proof process. This model has been used to specify prey agents behaviour inside a prey/predator system [16]. The resulting GDT contains sixteen nodes. This system have been also verified using the produced GDT. As shown before, this model can be seen as a tool for agents design. That's why we are going to develop an interpreter which can directly simulate the behaviour of agents from their GDT. The idea is to obtain, as in TAEMS, a method for fast prototyping with validation in parallel.

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